# Multiconfigurational methods: past, present and the road ahead



### Stefan Knecht





### Content of the lectures

- Concepts
- Past
  - CASCI
  - MCSCF/CASSCF
  - MR Dynamical Correlation Approaches

Present: "Sample" CI coefficients using an advanced optimisation algorithm

• Density Matrix Renormalisation Group

Road ahead: Beyond "classical quantum chemistry"

• Quantum Chemistry on Quantum Computers

# The Present ...



### Second example: the Cr<sub>2</sub> puzzle resolved

H. R. Larsson, H. Zhai, C. J. Umrigar, G. K.-L. Chan, JACS, 144, 15932 (2022)

### Cr<sub>2</sub> potential energy curse...



FIG. 1. Some of the simulated potential energy curves (PECs) of the chromium dimer that are available in the literature. The PECs are labeled by the study's year. The red curve marks the PEC from this work. The inset shows selected PECs from 2011 onwards. (List of references in SI).

Short and weak bond with a narrow minimum around 1.68 Å & extended shelf at around 2.5 Å.  $\rightarrow$  Cr 4s and 3d AOs different in size, with the minimum corresponding mostly to 3d orbital interactions and the shelf to 4s orbital interactions.

### A unique bonding and its consequences

- Complex electronic structure arises from interplay of two types of electron correlation • type | — "static correlation": spin-coupling of the 12 valence electrons (3d + 4s shells) energy-driven degeneracy

  - type II "spatial correlation / dynamic correlation": need for a large basis to capture excitations involving non-valence orbitals

by exciting to higher lying orbitals

### **Challenge:** $\bullet$

"The problem is computationally challenging because **both the static and dynamic correlation** must be computed sufficiently well even for a <u>qualitatively</u> reasonable description."

- overlap-driven degeneracy
- -> formation of 3d-3d bonds requires the 3p electrons to move out of the same spatial region



### Computational approach I

 $\Delta E = \Delta E_{\text{PDZ}}(\text{``exact''}) + \Delta E_{\text{CBS}}(\text{MPS-REPT2}) - \Delta E_{\text{DZ}}(\text{MPS-REPT2})$ 

- Static electron correlation: CAS(28e,76o)/cc-pVDZ-DK
  - Heat-Bath CI (selective CI)
  - DMRG (with huge bond dimension!)
- Dynamic electron correlation: CAS(12,12)/MRPT2
  - PT2 correlation of 3s and 3p inner-valence shells
  - all secondary shells considered

Basis set: MRPT2 with cc-pvNZ-DK (N=2,3,4,5) with extrapolation to CBS limit

### Computational approach II



### New state-of-the-art



$$\epsilon_{\text{REPT}} = |\Delta E(\text{MPS-REPT2}) - \Delta E[\text{UCCSD}($$

# Thinking outside the box **Standard CI approach**

• CI-type diagonalization for a **preselected** set of many-particle basis states  $k_1, k_2, \dots, k_L$ 

• Could we do better or, say, something else?

### $|\Psi\rangle = \sum c_{k_1,k_2,\ldots,k_L} |k_1\rangle \otimes |k_2\rangle \otimes \ldots \otimes |k_L\rangle$

# Thinking outside the box **Standard CI approach**

• CI-type diagonalization for a **preselected** set of many-particle basis states  $k_1, k_2, ..., k_L$ 

### • Could we do better or, say, something else?

- $\rightarrow$  selected CI approaches  $\rightarrow$  difference dedicated CI,  $\rightarrow$  many-body expansion FCI  $\rightarrow$  ...
- find "best" many-particle basis set based on correlations among orbitals  $\rightarrow$  **DMRG**

# $|\Psi\rangle = \sum c_{k_1,k_2,\ldots,k_L} |k_1\rangle \otimes |k_2\rangle \otimes \ldots \otimes |k_L\rangle$

• select many-particle basis sets / configurations / etc. based on some energy/weighting criteria:

# Thinking outside the box **Standard CI approach**

• CI-type diagonalization for a **preselected** set of many-particle basis states  $|\Psi\rangle = \sum c_{k_1,k_2,\ldots,k_L} |k_1\rangle \otimes |k_2\rangle \otimes \ldots \otimes |k_L\rangle$  $k_1, k_2, \dots, k_L$ 

• **Determine** CI coefficients from correlations among orbitals

$$|\Psi\rangle = \sum_{\substack{k_1,k_2,\ldots,k_L}} c_{k_1,k_2,\ldots,k_L}$$

• Local space  $k_l$  of *l*-th spatial orbital is of dimension d = 4:

### DMRG

- $_{k_L} | k_1 \rangle \otimes | k_2 \rangle \otimes \ldots \otimes | k_L \rangle$
- $k_{l} = \{ | \uparrow \downarrow \rangle | \uparrow \rangle, | \downarrow \rangle, | 0 \rangle \}$

From a CI wave function to Matrix Product States (MPS): Optimising an MPS with the DMRG algorithm

### Optimising an MPS wave function with the DMRG algorithm

- Optimisation algorithm
- Parameters that determine DMRG accuracy

# matrix product states, Annals of Physics, 326 (2011) 96–192.

U. Schollwöck, The density-matrix renormalization group in the age of



### Some reviews on about 20 years of DMRG in quantum chemistry

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### Intermission: singular value decomposition

• Singular value decomposition (SVD) of a matrix  $\mathbf{M}(n_a \times n_b)$ 

yields:

 $n_a$ 

- Left-singular matrix  $\mathbf{U}(n_a \times \min(n_a, n_b))$  with  $\mathbf{U}^{\dagger}\mathbf{U} = 1$
- Right-singular matrix  $\mathbf{V}(\min(n_a, n_b) \times n_b)$  with  $\mathbf{V}^{\dagger}\mathbf{V} = 1$
- Diagonal singular value matrix  $S(\min(n_a, n_b) \times \min(n_a, n_b))$  with r nonzero singular values  $\rightarrow r$  is the (Schmidt) rank of M  $n_h$  $n_a$  $n_a$  $n_b$



### Remember?



### From a CI to an MPS parametrisation I

• Successive application of SVD to CI tensor  $\rightarrow$  MPS wave function







- S and V<sup>†</sup> multiplied and reshaped into coefficient tensor  $C_{a_1,(k_2,...,k_L)}$
- $r_1 \leq d$
- collection of d(=4) row vectors  $A^k$

$$k_L = c_{k_1,k_2,\ldots,k_L}$$

$$U_{k_1,a_1}S_{a_1,a_1}(V^{\dagger})_{a_1,(k_2,\ldots,k_L)}$$

$$A_{1,a_1}^{k_1} c_{a_1,(k_2,...,k_L)}$$

$$^{k_1}$$
 with entries  $A_{1,a_1}^{k_1} = U_{k_1,a_1}$ 

### From a CI to an MPS parametrisation III • Reshape coefficient tensor $c_{a_1,(k_2,\ldots,k_L)}$ into a $r_1 d \times d^{L-2}$ matrix $\Gamma$



- S and V<sup>†</sup> multiplied and reshaped into coefficient tensor  $C_{a_1,(k_2,\ldots,k_L)}$
- $r_2 \leq r_1 d \leq d^2$
- collection of d matrices  $A^{k_2}$  with en

 $\underset{=}{\overset{'1}{\sum}} \sum_{n=1}^{\prime 2} A_{1,a_1}^{k_1} U_{(a_1k_2),a_2} S_{a_2,a_2} (V^{\dagger})_{a_2,(k_3,\ldots,k_L)}$ 

$$\sum_{i=1}^{r_2} A_{1,a_1}^{k_1} A_{a_1,a_2}^{k_2} \Gamma_{(a_2k_3),(k_4,\ldots,k_L)}$$

tries 
$$A_{a_1,a_2}^{k_2} = U_{(a_1k_2),a_2}$$

### From a CI to an MPS parametrisation IV

### • Continue with SVDs until last site which then gives

$$c_{k_1,k_2,\dots,k_L} = \sum_{a_1,a_2,\dots,a_L-1} A_{1,a_1}^{k_1} A_{a_1,a_2}^{k_2} \cdots A_{a_{L-2,L-1}}^{k_{L-1}} A_{a_{L-1},1}^{k_L}$$
$$\equiv A^{k_1} A^{k_2} \cdots A^{k_{L-1}} A^{k_L}$$

### with

- interpretation of sums as matrix-matrix multiplications
- first and last matrices are row- and column vectors!
- CI wave function rewritten as MPS wave function:

$$\left|\Psi\right\rangle = \sum_{k} c_{k} \left|k\right\rangle = \sum_{k_{1},k_{2},\ldots,k_{L}} A^{k_{1}}A^{k_{2}}\cdots A^{k_{L-1}}A^{k_{L}} \left|k\right\rangle$$

### From a CI to an MPS parametrisation V





### Properties of the MPS I

- Matrix dimensions grow exponentially up to  $dim(d^{L/2-1} \times d^{L/2})$  if no truncation occurs, i.e., all singular values are kept
- From  $U^{\dagger}U = I$  follows that all matrices  $\{A^{k_l}\}$  are left-normalised
- MPS built from left-normalised matrices is called left-canonical
- For any lattice bipartition at site l, the states on sites  $1, \ldots, l$

$$|a_l\rangle_{\mathscr{L}} = \sum_{k_1, k_2, \dots, k_l} \left(A^{k_1} \cdots A^{k_l}\right)_{1, a_l} |k_1, \dots, k_l\rangle$$

span a left subsystem  $\mathscr{L}$  and form an orthonormal basis



### Properties of the MPS II

- Starting SVD on coefficient tensor from right-hand side  $I'_{(k_1,k_2,...,k_{L-1})}$ 
  - yields right-normalised matrices { **)** *E*
- MPS built from right-normalised matrices is called right-canonical
- For any lattice bipartition at site l + 1, the states on sites  $l + 1 \dots L$  $|a_{l+1}\rangle_{\mathcal{R}} = \sum \left(B^{k_{l+1}}\cdots B^{k_L}\right)_{a_{l+1}}|k_{l+1},\ldots,k_L\rangle$  $k_{l+1}, k_{l+2}, \dots, k_L$

span a right subsystem  $\mathscr{R}$  and form an orthonormal basis

$$_{1}),k_{L} = c_{k_{1},k_{2},...,k_{L}}$$

$$B^{k_{l}} \{ as V^{\dagger}V = I \}$$

$$B^{k_{l}}B^{k_{l}\dagger} = I$$

### Gauge freedom and mixed-canonical form

- MPS representations are not unique  $\leftrightarrow$  existence of a gauge degree of freedom
- Consider two adjacent matrices  $M^{k_l}$  and  $M^{k_{l+1}}$  of shared column/row dimension D and a square invertible matrix X ( $D \times D$ )
- Invariance of MPS immediately follows from

$$M^{k_l} \rightarrow M^{k_l} X;$$

since

$$M^{k_l} \underbrace{XX^{-1}}_{M^{k_{l+1}}} M^{k_{l+1}} = M^{k_l} \cdot M^{k_{l+1}}$$

$$M^{k_{l+1}} \to X^{-1}M^{k_{l+1}}$$



### Mixed-canonical MPS representation

• Gauge freedom allows to write an MPS in mixed canonical form at sites  $\{l, l+1\}$ 

$$\left|\Psi\right\rangle = \sum_{k} A^{k_{1}} \cdots A^{k_{l-1}} M^{k_{l}k_{l+1}} B^{k_{l+2}} \cdots B^{k_{L}} \left|k\right\rangle$$

by starting from a general MPS wave function

$$|\Psi\rangle = \sum_{k} M^{k_1} M^{k_2} \cdots M^{k_L} |k\rangle$$

and the **two-site MPS tensor**  $M^{k_l k_{l+1}}$  reading as

$$M^{k_l k_{l+1}} \equiv M^{k_l k_{l+1}}_{a_{l-1}, a_{l+1}} = \sum_{a_l} M^{k_l}_{a_{l-1}, a_l} M^{k_{l+1}}_{a_l, a_{l+1}}$$



### Matrix product operators I

• MPS concept applied to operators  $\rightarrow$  matrix product operators (MPOs)



- N-electron operator  $\widehat{\mathcal{M}}$  in MPO form  $\widehat{\mathcal{W}} = \sum W$ 
  - $kk' b_1, ..., b_{L-1}$

 $= \sum W^{k_1k_1'}W^{k_2k_2}$ kk'

 $\equiv \sum w_{kk'} \left| k \right\rangle \left\langle k' \right|$ kk'



$$V_{1,b_1}^{k_1k_1'}W_{b_1,b_2}^{k_2k_2'}\cdots W_{b_{L-1},1}^{k_Lk_L'} \left| \boldsymbol{k} \right\rangle \left\langle \boldsymbol{k'} \right|$$

$$k'_2 \cdots W^{k_L k'_L} | \mathbf{k} \rangle \langle \mathbf{k'} |$$

### Matrix product operators II

the local site indices  $k_l k'_l$ 

This allows us to write the equation on previous slide as

 $b_1, ..., b_{L-1}$ • Note: the entries of  $\{W_{b_{l-1},b_l}\}$  matrices comprise the elementary, *local* operators acting on the *l*-th orbital, e.g.,  $\tilde{a}_{\uparrow}^{\dagger} = |\uparrow\downarrow\rangle$ 

For efficiency, rearrange summations such that the contraction proceeds first over

$$\widehat{W_{b_{l-1},b_l}^l} = \sum_{\substack{k_l k_l'}} W_{b_{l-1},b_l}^{k_l k_l'} \left| k_l \right\rangle \left\langle k_l' \right|$$

$$\widehat{\mathcal{W}} = \sum \widehat{W_{1,b_1}^1} \cdots \widehat{W_{b_{l-1},b_l}^l} \cdots \widehat{W_{b_{L-1},1}^L}$$

$$\rangle \langle \downarrow | + | \uparrow \rangle \langle 0 |$$







- Prerequisite: initialise suitable (valid) trial MPS wave function  $|\Psi\rangle$ 
  - choices: random guess, encode HF determinant, CI-DEAS, "old MPS" ...
  - assume normalisation, i.e.,  $\langle \Psi | \Psi \rangle = 1$





- Ansatz for variational MPS optimization: extremize the Lagrangian
  - $\mathscr{L} = \langle \Psi | \hat{H} | \Psi$

with the two-site  $\{M^{k_l k_{l+1}}\}$  matrices as optimization parameters

- Optimize at each step of a "sweep" entries of site matrices of <u>two orbitals</u> ("two-site DMRG") while keeping all the others fixed
- Sweep through all sites multiple times until energy converges

$$\Psi \rangle - \lambda \left( \langle \Psi | \Psi \rangle - 1 \right)$$

### Variational MPS optimisation III

 $M^{k_l,k_{l+1}}$  $\frac{\partial}{\partial M^{k_{l},k_{l+1}*}} \left( \left\langle \Psi \right| \hat{H} \right| \Psi$ 

which then yields

 $\sum_{\substack{a'_{l-1}a'_{l}k'_{l+1}}} \sum_{\substack{k'_{l+1}k'_{l+1}}} L^{b_{l-1}}_{a_{l-1},a'_{l-1}} V$  $b_{l-1}b_{l+1}$ 

and ket MPS starting from left (right) up to sites l - 1(l + 1)

• At sites  $\{l, l+1\}$ , take derivative of  $\mathscr{L}$  with respect to complex conjugate of

$$\Psi \rangle - \lambda \left[ \langle \Psi | \Psi \rangle - 1 \right] \right) = 0$$

$$W_{b_{l-1},b_{l+1}}^{k_{l}k_{l+1},k_{l}'k_{l+1}'}R_{a_{l+1}',a_{l+1}}^{b_{l+1}}M_{a_{l-1}',a_{l+1}'}^{k_{l}'k_{l+1}'} = \lambda \sum_{a_{l-1}'} \Psi_{a_{l-1}',a_{l-1}}^{A}$$

$$imes M^{k'_l k'_{l+1}}_{a'_{l-1},a'_{l+1}}$$

 $\times \Psi^{B}_{a'}$  $a'_{l+1}, a_{l+1}$ 

• L and R: left and right boundaries obtained by contracting the MPO with the bra



# Variational MPS optimisation IV







### Variational MPS optimisation V

$$\sum_{\substack{a_{l-1}'a_{l}' \\ b_{l-1}b_{l+1}}} \sum_{\substack{k_{l}'k_{l+1}' \\ b_{l-1}b_{l+1}}} L_{a_{l-1},a_{l-1}'}^{b_{l-1}} W_{b_{l-1},b_{l+1}}^{k_{l}k_{l+1}'} R_{a_{l+1}',a_{l+1}}^{b_{l+1}} M_{a_{l-1}',a_{l+1}'}^{k_{l}'k_{l+1}'} = \lambda M_{a_{l-1}',a_{l+1}'}^{k_{l}'k_{l+1}'}$$
**if** MPS is a canonical MPS!

R

• Requires the initial MPS to be right-normalized!

• NB: Simplify generalized eigenvalue problem to a standard eigenvalue problem



### Variational MPS optimisation VI

- Recast last equation into a matrix eigenvalue equation  $\mathcal{H}v - \lambda v = 0$ 
  - by defining a local Hamiltonian matrix  $\mathcal{H}$  at sites  $\{l, l+1\}$
  - and a vector v

 $V_{k'_{l}k'_{l+1}a'_{l-1}a'_$ 

# $H_{(k_{l}k_{l+1}a_{l-1}a_{l+1}),(k_{l}'k_{l+1}'a_{l-1}'a_{l+1}')} = \sum L_{a_{l-1},a_{l-1}'}^{b_{l-1}} W_{b_{l-1},b_{l+1}}^{k_{l}k_{l+1},k_{l}'k_{l+1}'} R_{a_{l+1}',a_{l+1}}^{b_{l+1}}$ $b_{l-1}, b_{l+1}$

$$a_{l+1}' = M_{a_{l-1}',a_{l+1}'}^{k_l'k_{l+1}'}$$

• Solving EV problem —> eigenvalue  $\lambda^0$  and corresponding eigenvector  $v_{k'_lk'_{l+1}a'_{l-1}a'_{l+1}}^0$ 



### Variational MPS optimisation VII

- Reshape  $v_{k'_lk'_{l+1}a'_{l-1}a'_{l+1}}^0$  back to  $M_{a'_{l-1}a'_{l+1}}^{k'_lk'_{l+1}}$
- $M_{a'_{l-1},a'_{l-1}}^{k'_{l+1}}$  is subsequently subject to a left- or right-normalisation (SVD!)

$$M_{a'_{l-1},a'_{l+1}}^{k'_{l}k'_{l+1}} = M_{(k'_{l},a'_{l-1})(k'_{l+1},a'_{l+1})} = U_{(k'_{l},a'_{l-1})s_{l}}S_{s_{l}s_{l}}V_{s_{l}(a'_{l+1},k'_{l+1})}$$

- By discarding the 3m smallest singular values in  $S_{S_1S_1}$  to obtain  $S_{a_1a_1}$  we achieve the desired reduction in bond dimensionality!
- The maximum (fixed) number *m* of retained singular values is usually called number of renormalized block states

### Variational MPS optimisation VIII

• Discarding 3m smallest singular values corresponds to discarding the last 3m columns (rows) of U(V) such that

$$A_{a'_{l-1},a'_{l}}^{k'_{l}} \equiv U_{(k'_{l},a'_{l-1})}$$

$$I - S_l =$$

Energy calculated as a function of the truncation error  $\epsilon$  $\bullet$ 

$$\epsilon = \sum_{\substack{s_l = m+1}}^{4m} S_{s_l s_l} = ||\Psi_{16m^2} - \Psi_{4m^2}||$$

can be employed to obtain an error estimate through extrapolation



### Variational MPS optimisation IX

• Moving from sites  $\{l, l + 1\}$  to sites optimization step

### • Moving from sites $\{l, l+1\}$ to sites $\{l+1, l+2\}$ then completes the local
# Scaling of variational MPS optimisation

- Scaling is dominated by cost of contracting the operator with the MPS on one site and is proportional to the number of non-zero elements in the MPO matrices  $\{W\}$ 
  - in a naïve MPO ansatz this step scales as  $\mathcal{O}(L^5)$
  - in an optimized code scaling reduces to  $\mathcal{O}(L^4)$
- Further reduction through symmetry: U(1) and SU(2)
- SVD scales as  $\mathcal{O}(m^3)$  (but there are L of them in a sweep)
- Taking into account all operations a sweep scales  $\approx O(L^4 m^3)$



## Extrapolation

- Extrapolate E based on truncation error  $\epsilon$  for different values of m
- Example: ground-state calculation of F<sub>2</sub>



# $\ln\left(\frac{E_{\text{DMRG}} - E_{\text{FCI}}}{E_{\text{FCI}}}\right) = a \ln \epsilon + b$

# Determining factors of DMRG convergence

- Size L of the CAS
- Type of molecular orbitals (HF, NO's, localized orbitals, ...)
- MPS guess for the right subsystem (initial sweep)
- Ordering of orbitals (exploit quantum information / graph theory)
- Number of renormalized block states m

various runs until results converge!

One should never calculate results for just a single *m*, but increase it in

# Properties of DMRG

- Variational
- Size-consistent
- (approximate) FCI for a CAS
- Polynomial scaling (  $\approx L^4 m^3$ )
- MPS wave function
- For large *m* invariant wrt orbital rotations

- Variational
- Size-consistent
- FCI for a CAS
- Factorial scaling
- Linearly parametrised wave function
- Invariant wrt orbital rotations



# (Incomplete) List of DMRG codes for QC

- Budapest-DMRG code (Matlab, no source code available)
- MOLMPS (C++, ?)
- (Stack)Block and Block2 (C++, source code available)
- CheMPS2 (C++, Fortran, source code available)
- QCMaquis (C++, Fortran, source code available)
- •

# Other classical methods for large CAS

- FCI-Quantum Monte Carlo aka FCIQMC
- Heat-Bath CI (aka SHCI)
- selective CI / CIPSI-like approaches
- v(ariational) 2RDM

## -> Extensions to treat dynamical electron correlation available!

# The road ahead ...



## Thinking outside the box II **Standard CI approach**

• CI-type diagonalization for a **preselected** set of many-particle basis states

$$|\Psi\rangle = \sum_{k_1,k_2,\ldots,k_L} c_{k_1,k_2,\ldots,k_L}$$

• **Determine** CI coefficients from correlations among orbitals

$$\left|\Psi\right\rangle = \sum_{k_1,k_2,\ldots,k_L} c_{k_1,k_2,\ldots,k_L} \left|k_1\right\rangle \otimes \left|k_2\right\rangle \otimes \ldots \otimes \left|k_L\right\rangle$$

 $\langle \Psi | \hat{H}_e | \Psi \rangle =$ 

using an entangled set of qubits

• "Learn" the energy

## $k_L | k_1 \rangle \otimes | k_2 \rangle \otimes \ldots \otimes | k_L \rangle$ IRG

## Quantum Computing

$$= \sum_{k} c_{k} \langle \Psi \mid P_{k} \mid \Psi \rangle$$

 $|\Psi\rangle = \mathbf{U}(\vec{\theta}) |k_1\rangle \otimes |k_2\rangle \otimes \ldots \otimes |k_L\rangle$ 

Quantum Chemistry on a Quantum Computer: Concepts and Challenges

# Some important references

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Yuri Manin 1980



Richard Feynman 1982



Yuri Manin 1980

## Simulating some quantum mechanical effects on a classical computer is unfeasible



Richard Feynman 1982

#### Use a quantum one!



- Classical bit can be **either** in state  $|0\rangle$  or state  $|1\rangle$ 
  - Qubit can be in a superposition of both states

$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle \equiv \alpha \begin{bmatrix} 1\\0 \end{bmatrix} + \beta \begin{bmatrix} 0\\1 \end{bmatrix}$$





From bits to ... --> qubits

$$|\psi\rangle = |0\rangle \xrightarrow{\text{superposition}} |\psi\rangle = \alpha |0\rangle + \beta |1$$

- Classical bit can be **either** in state  $|0\rangle$  or state  $|1\rangle$ 
  - Qubit can be in a **superposition** of both states  $|\psi\rangle = \alpha |0\rangle + \beta |1\rangle \equiv \alpha \begin{bmatrix} 1\\ 0 \end{bmatrix} + \beta \begin{bmatrix} 0\\ 1 \end{bmatrix}$







From bits to	qubits	• Mea

Р

$$\begin{vmatrix} \mathbf{\psi} \\ \mathbf{\psi} \\ = \begin{vmatrix} 0 \\ \mathbf{\psi} \end{vmatrix} = \alpha \begin{vmatrix} 0 \\ \mathbf{\psi} \\ \mathbf{\psi} \end{vmatrix} = \alpha \begin{vmatrix} 0 \\ \mathbf{\psi} \\ \mathbf{\psi} \end{vmatrix} = \alpha \begin{vmatrix} 0 \\ \mathbf{\psi} \\ \mathbf{\psi} \end{vmatrix} = \alpha \begin{vmatrix} 0 \\ \mathbf{\psi} \\ \mathbf{\psi} \\ \mathbf{\psi} \end{vmatrix}$$

- Classical bit can be **either** in state  $|0\rangle$  or state  $|1\rangle$ 
  - Qubit can be in a superposition of both states

$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle \equiv \alpha \begin{bmatrix} 1\\0 \end{bmatrix} + \beta \begin{bmatrix} 0\\1 \end{bmatrix}$$

asuring the state of the qubit with probability P

$$\left( \left| \psi \right\rangle = \left| 0 \right\rangle \right) = \gamma^2 \text{ and } P\left( \left| \psi \right\rangle = \left| 1 \right\rangle \right) = \delta^2$$







Arbitrary state of its qubits

 $\rightarrow |\Psi\rangle = \sum_{\vec{f}} \alpha_{\vec{f}} |f_1 f_2 \dots f_M\rangle$ 



Arbitrary state of its qubits

 $\rightarrow |\Psi\rangle = \sum_{\vec{r}} \alpha_{\vec{f}} |f_1 f_2 \dots f_M\rangle \quad \rightarrow$ 

A universal quantum computer can solve problems beyond quantum simulation (e.g. factorisation)





Arbitrary state of its qubits

 $\rightarrow |\Psi\rangle = \sum_{\vec{J}} \alpha_{\vec{f}} |f_1 f_2 \dots f_M\rangle \quad \rightarrow$ 

A universal quantum computer can solve problems beyond quantum simulation (e.g. factorisation)



Quantum 5, 433 (2021)





Arbitrary state of its qubits

 $\rightarrow |\Psi\rangle = \sum_{\vec{J}} \alpha_{\vec{f}} |f_1 f_2 \dots f_M\rangle \quad \rightarrow$ 

A universal quantum computer can solve problems beyond quantum simulation (e.g. factorisation)



Quantum 5, 433 (2021)







Arbitrary state of its qubits

 $\rightarrow |\Psi\rangle = \sum_{\vec{f}} \alpha_{\vec{f}} |f_1 f_2 \dots f_M\rangle \quad \rightarrow$ 

We are in the era of the Noisy Intermediate-Scale Quantum computers: soon useful for simulation!

A universal quantum computer can solve problems beyond quantum simulation (e.g. factorisation)

![](_page_55_Figure_6.jpeg)

Quantum 5, 433 (2021)

![](_page_55_Picture_8.jpeg)

#### Quantum computer technologies Select Players าร eds and Google IBM Q riqetti ge standard es. Among OQC IQM so has a QuTech ogenic q|c|1ence times; ○本源量子 Origin Quantum ect antum Circuits, In well h gate **O**AQT IONQ a herence genic lons are $\bigcirc$ nt. es/ QUANTINUUM connectivity ers hard to Control in the second s Universal Quantum a-high n charges τy. gate ng fidelities. $\otimes$ Ψ PsiQuantum cuums rall footprint. $X \land N \land D U$ ng CMOS noton loss; es its own QUANTUM naturally challenges. ce times. QuEra> ۱d ColdQuanta onnectivity, COMPUTING INC. 2Q. External 쥖 red. atom a-high Computing aling PASQAL isting Silicon nology. (intel) Quantum Computing and diraq ogenics. R d gates toence times.

Qubit Type	Pros/Cons
Superconducting	Pros: High gate speeds fidelities. Can leverage lithographic processes. first qubit modalities so head start. Cons: Requires cryoge cooling; short coherence microwave interconnect frequencies still not we understood.
Trapped lons	<b>Pros</b> : Extremely high g fidelities and long cohe times. Extreme cryoge cooling not required. lo perfect and consistent. <b>Cons</b> : Slow gate times operations and low con between qubits. Lasers align and scale. Ultra-h vacuum required. lon c may restrict scalability.
Photonics	<b>Pros</b> : Extremely fast gas speeds and promising to No cryogenics or vacuu required. Small overall Can leverage existing ( fabs. <b>Cons</b> : Noise from phot
	each program requires chip. Photons don't nat interact so 2Q gate cha
Neutral Atoms	<b>Pros</b> : Long coherence Atoms are perfect and consistent. Strong conr including more than 2C cryogenics not required <b>Cons</b> : Requires ultra-h vacuums. Laser scalin
Silicon Spin/Quantum Dots	challenging. <b>Pros</b> : Leverages existing semiconductor technology Strong gate fidelities and speeds. <b>Cons</b> : Requires cryoger Only a few entangled go date with low coherence Interference/cross-talkor challenges.

QUANTUM MOTION

QUANTUM BRILLIANCE

Source: Quantum Computing Modalities -A Qubit Primer Revisited -The Quantum Leap (quantum tech.blog)

![](_page_56_Picture_3.jpeg)

![](_page_57_Picture_0.jpeg)

## Electronic wave-function representation

Fermion-to-qubit mapping

#### QUANTUM STATE

![](_page_57_Picture_4.jpeg)

### Quantum Computing for Quantum Chemistry in a Nutshell

The state of the quantum processor mathematically represents the state of the molecule

The goal is to find the state of the molecule for which the energy is minimal

The energy of the molecule needs to be measured

![](_page_57_Picture_9.jpeg)

![](_page_58_Picture_0.jpeg)

## Electronic wave-function representation

Fermion-to-qubit mapping

![](_page_58_Picture_3.jpeg)

## Challenges in near-term QC

- 1. Measurement stage is time-consuming
- 2. Hilbert space is a big space
- 3. Qubits are a scarce resource
- 4. Noise biases the results
- 5. Resource efficient/aware representation

![](_page_58_Picture_10.jpeg)

![](_page_59_Picture_0.jpeg)

## Challenges in near-term QC

- 1. Measurement stage is time-consuming
- 2. Hilbert space is a big space
- 3. Qubits are a scarce resource
- 4. Noise biases the results
- 5. Resource efficient/aware representation

![](_page_59_Figure_7.jpeg)

![](_page_60_Picture_0.jpeg)

## Electronic wave-function representation

Fermion-to-qubit mapping

#### QUANTUM STATE

![](_page_60_Picture_4.jpeg)

## Challenges in near-term QC

- 1. Measurement stage is time-consuming
- 2. Hilbert space is a big space
- 3. Qubits are a scarce resource
- 4. Noise biases the results
- 5. Resource efficient/aware representation

![](_page_60_Picture_11.jpeg)

![](_page_61_Picture_0.jpeg)

## Challenges in near-term QC

- 1. Measurement stage is time-consuming
- 2. Hilbert space is a big space
- 3. Qubits are a scarce resource
  - 4. Noise biases the results
  - 5. Resource efficient/aware representation

n

![](_page_62_Picture_0.jpeg)

Electronic wave-function representation

Fermion-to-qubit mapping

#### QUANTUM STATE

![](_page_62_Picture_4.jpeg)

## Challenges in near-term QC

- 1. Measurement stage is time-consuming
- 2. Hilbert space is a big space
- 3. Qubits are a scarce resource
- 4. Noise biases the results

5. Resource efficient/aware representation

![](_page_62_Picture_11.jpeg)

# Mapping "the problem" from fermion to qubit space

## Mapping the problem from fermion to qubit space General considerations

N-electron wave function

$$\left|\Psi\right\rangle = \sum_{\zeta} c_{\zeta} \left|\Phi_{\zeta}\right\rangle$$

Second-quantized Hamiltonian

$$\hat{H}_{e} = \sum_{p,q} h_{pq} \, \hat{E}_{pq} + \frac{1}{2} \sum_{p,q,r,s} (pq \,|\, rs) \, \left( \hat{E}_{pq} \hat{E}_{rs} - \delta_{qr} \right)$$

• Fermions are indistinguishable particles, qubits are distinguishable

$$\{\hat{a}_{p},\hat{a}_{q}\}=0$$
  $\{\hat{a}_{p}^{\dagger},\hat{a}_{q}^{\dagger}\}=0$   $\{\hat{a}_{p},\hat{a}_{q}^{\dagger}\}=\delta_{pq}$ 

![](_page_64_Figure_7.jpeg)

-> we need to account for anti-commutation of fermionic operators in the map!

![](_page_64_Picture_9.jpeg)

## Mapping the problem from fermion to qubit space *N*-electron wave function

$$\left|\Psi\right\rangle = \sum_{\mathbf{f}} c_{\mathbf{f}} \left|f_{M-1} \dots f_{1} f_{0}\right\rangle$$

Exact for infinite *M*!

![](_page_65_Picture_3.jpeg)

## Mapping the problem from fermion to qubit space *N*-electron wave function

$$\left|\Psi\right\rangle = \sum_{\mathbf{f}} c_{\mathbf{f}} \left|f_{M-1} \dots f_{1} f_{0}\right\rangle$$

Trivial interaction of fermionic creation/annihilation operators on ON

$$\hat{a}_{p}^{\dagger} \left| f_{M-1} \dots f_{p+1} \ 0 \ f_{p-1} \dots f_{0} \right\rangle = (-1)^{\sum_{s=0}^{p-1} f_{s}} \left| f_{M-1} \dots f_{p+1} \ 1 \ f_{p-1} \dots f_{0} \right\rangle$$

$$\hat{a}_{p}^{\dagger} \left| f_{M-1} \dots f_{p+1} \ 1 \ f_{p-1} \dots f_{0} \right\rangle = 0$$

$$\hat{a}_{p} \left| f_{M-1} \dots f_{p+1} \ 1 \ f_{p-1} \dots f_{0} \right\rangle = (-1)^{\sum_{s=0}^{p-1} f_{s}} \left| f_{M-1} \dots f_{p+1} \ 0 \ f_{p-1} \dots f_{0} \right\rangle$$
$$\hat{a}_{p} \left| f_{M-1} \dots f_{p+1} \ 0 \ f_{p-1} \dots f_{0} \right\rangle = 0$$

### Exact for infinite *M*!

![](_page_66_Picture_7.jpeg)

## Mapping the problem from fermion to qubit space *N*-electron wave function

Trivial translation of fermionic ON vector basis to qubit ON vector basis

$$|f_{M-1}...f_1f_0\rangle \longrightarrow |q_{M-1}\rangle...\otimes |q_1\rangle \otimes |q_0\rangle \equiv |q_{M-1}...q_1q_0\rangle$$

What about creation/annihilation operators for qubits?  $\bullet$ 

![](_page_67_Picture_4.jpeg)

Use a simple recipe for one-qubit creation and annihilation operators?

$$\hat{Q}^{+} \left| 1 \right\rangle = 0$$
  $\hat{Q}^{+} \left| 0 \right\rangle = \left| 1 \right\rangle$   $\hat{Q}^{-} \left| 1 \right\rangle = \left| 0 \right\rangle$   $\hat{Q}^{-} \left| 0 \right\rangle = 0$ 

![](_page_68_Picture_4.jpeg)

Use a simple recipe for one-qubit creation and annihilation operators? 

$$\hat{Q}^{+} |1\rangle = 0$$
  $\hat{Q}^{+} |0\rangle = |1\rangle$   $\hat{Q}^{-} |1\rangle = |0\rangle$   $\hat{Q}^{-} |0\rangle = 0$ 

Do not obey the fermionic anti-commutation relations!

![](_page_69_Picture_4.jpeg)

![](_page_69_Picture_6.jpeg)

Use a simple recipe for one-qubit creation and annihilation operators? 

$$\hat{Q}^{+} |1\rangle = 0$$
  $\hat{Q}^{+} |0\rangle = |1\rangle$   $\hat{Q}^{-} |1\rangle = |0\rangle$   $\hat{Q}^{-} |0\rangle = 0$ 

Do not obey the fermionic anti-commutation relations!

Better ansatz? Form operators based on Pauli matrices!

$$\hat{Q}^{+} = |1\rangle\langle 0| \equiv \frac{1}{2}\left(\sigma_{x} - i\sigma_{y}\right) = \frac{1}{2}\left(X - iY\right)$$
$$\hat{Q}^{-} = |0\rangle\langle 1| \equiv \frac{1}{2}\left(\sigma_{x} + i\sigma_{y}\right) = \frac{1}{2}\left(X + iY\right)$$

$$\hat{Q}^{+} = |1\rangle \langle 0| \equiv \frac{1}{2} \left(\sigma_{x} - i\sigma_{y}\right) = \frac{1}{2} \left(X - iY\right)$$
$$\hat{Q}^{-} = |0\rangle \langle 1| \equiv \frac{1}{2} \left(\sigma_{x} + i\sigma_{y}\right) = \frac{1}{2} \left(X + iY\right)$$

![](_page_70_Picture_7.jpeg)

![](_page_70_Picture_9.jpeg)

• Why are Pauli matrices a suitable choice?

![](_page_71_Figure_4.jpeg)

![](_page_71_Picture_5.jpeg)
• Why are Pauli matrices a suitable choice? Single-qubit quantum gates are  $2 \times 2$  unitary matrices

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \qquad \sigma_y =$$

$$\begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \qquad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$



Jordan-Wigner mapping

- Why are Pauli matrices a suitable choice? Single-qubit quantum gates are  $2 \times 2$  unitary matrices
  - $= \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \qquad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$

  - -> Pauli matrices are  $2 \times 2$  Hermitian (unitary) matrices −> Pauli matrices are involutory:  $\sigma_p^2 = \mathbf{1} \forall p \in [x, y, z]$ -> Pauli matrices anti-commute:  $\{\sigma_p, \sigma_q\} = 2I\delta_{pq} \quad \forall p, q \in [x, y, z]$

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \qquad \sigma_y =$$

# Mapping the problem from fermion to qubit space



- Because of anti-commutation among Paulis,  $\hat{Q}^{\pm}$  and  $\sigma_{_7}$  anti-commute
- Represent action of fermionic operators  $\hat{a}_p^{(\dagger)}$  for index p by acting with  $\hat{Q}^{\pm}$  on qubit p and with  $\sigma_z$  on all qubits with index q < p and with the identity  $\mathbf{1}$  on the remaining qubits



Jordan-Wigner mapping

- Because of anti-commutation among Paulis,  $\hat{Q}^{\pm}$  and  $\sigma_{_7}$  anti-commute
- Represent action of fermionic operators  $\hat{a}_p^{(\dagger)}$  for index p by acting with  $\hat{Q}^{\pm}$  on qubit p and with  $\sigma_{z}$  on all qubits with index q < p and with the identity  $\mathbf{1}$  on the remaining qubits

$$\hat{a}_{p}^{\dagger} \equiv \mathbf{1}^{\otimes M - p - 1} \otimes \hat{Q}_{p}^{+} \otimes \left[\sigma_{z}^{\otimes p}\right] = \frac{1}{2} \left(X_{p} \otimes \left[\sigma_{z}^{\otimes p}\right] - iY_{p} \otimes \left[\sigma_{z}^{\otimes p}\right]\right)$$

# Mapping the problem from fermion to qubit space







Jordan-Wigner mapping

- Because of anti-commutation among Paulis,  $\hat{Q}^{\pm}$  and  $\sigma_{_{\! 7}}$  anti-commute
- Represent action of fermionic operators  $\hat{a}_p^{(\dagger)}$  for index p by acting with  $\hat{Q}^{\pm}$  on qubit p and with  $\sigma_{z}$  on all qubits with index q < p and with the identity 1 on the remaining qubits

$$\hat{a}_{p}^{\dagger} \equiv \mathbf{1}^{\otimes M-p-1} \otimes \hat{Q}_{p}^{+} \otimes \left[\sigma_{z}^{\otimes p}\right] = \frac{1}{2} \left(X_{p} \otimes \left[\sigma_{z}^{\otimes p}\right] - iY_{p} \otimes \left[\sigma_{z}^{\otimes p}\right]\right)$$
$$\hat{a}_{p} \equiv \mathbf{1}^{\otimes M-p-1} \otimes \hat{Q}_{p}^{-} \otimes \left[\sigma_{z}^{\otimes p}\right] = \frac{1}{2} \left(X_{p} \otimes \left[\sigma_{z}^{\otimes p}\right] + iY_{p} \otimes \left[\sigma_{z}^{\otimes p}\right]\right)$$

# Mapping the problem from fermion to qubit space









Major drawback of Jordan-Wigner mapping: k-locality!

of extra qubit operations which scale as  $\mathcal{O}(M)$ !

**Example:** how does it work for mapping one-electron terms

$$\hat{h} = \sum_{pq} h_{pq} \hat{a}_p^{\dagger} \hat{a}_q ?$$

- k-locality: each term in the Hamiltonian acts non-trivially on at most k qubits
- -> non-locality of "parity term  $[\sigma_z^{\otimes p}]$ " that appears in  $\hat{a}_p^{(\dagger)}$  introduces a number



fcac = np.zeros((2, 2), dtype=float) fcac[0][0] = 1.0; fcac[1][1] = 3.0 h2ac = np.zeros((2, 2, 2, 2), dtype=float) > one\_body\_ints = OneBodyElectronicIntegrals( ---> two\_body\_ints = TwoBodyElectronicIntegrals( --fermi\_ham = qc\_ham\_qiskit.second\_q\_ops()[0] fermi\_ham.set\_truncation(0) print("Fermionic Hamiltonian",fermi\_ham) Hq = qubit\_converter.convert(fermi\_ham) print("J-W Qubit Hamiltonian",Hq)

√ 0.4s

Fermionic Hamiltonian Fermionic Operator register length=4, number terms=5

 $1.0 * (+_0 -_0)$  $+ 3.0 * (+_1 -_1)$ + 1.0 \* ( +\_2 -\_2 ) + 3.0 \* ( +\_3 -\_3 ) + 0j J-W Qubit Hamiltonian 4.0 \* IIII - 0.5 \* IIIZ 1.5 \* IIZI - 0.5 \* IZII • 1.5 \* ZIII

> from qiskit\_nature.converters.second\_quantization import QubitConverter

qc\_ham\_qiskit = ElectronicEnergy([one\_body\_ints, two\_body\_ints])

qubit\_converter = QubitConverter(mapper=JordanWignerMapper())



> from qiskit\_nature.converters.second\_quantization import QubitConverter fcac = np.zeros((2, 2), dtype=float) fcac[0][0] = 1.0; fcac[1][1] = 3.0

h2ac = np.zeros((2, 2, 2, 2), dtype=float) > one\_body\_ints = OneBodyElectronicIntegrals( ---

> two\_body\_ints = TwoBodyElectronicIntegrals( --qc\_ham\_qiskit = ElectronicEnergy([one\_body\_ints, two\_body\_ints]) fermi\_ham = qc\_ham\_qiskit.second\_q\_ops()[0]

# Could we avoid to operate with $\sigma_7$ on all qubits q with

	Q
гепштоптс наштегоптан гегш.	TOUTC
register length=4, number	terms=
$1.0 * (+_00)$	
+ 3.0 * ( +_11 )	
+ 1.0 * ( +_22 )	
+ 3.0 * ( +_33 )	
+ 0j	
J-W Qubit Hamiltonian 4.0 >	* IIII
- 0.5 * IIIZ	
– 1.5 * IIZI	
- 0.5 * IZII	
– 1.5 * ZIII	

operator



• Alternative idea to occupation number basis -> parity basis:

Use qubit p to store the parity  $\mathscr{P}$  of all occupied orbitals up to p

$$\mathscr{P}_p = \mod$$

phase of -1

-> it suffices to only have  $\sigma_z$  acting on qubit p-1

$$\left(\sum_{s=0}^{p} f_{s}, 2\right)$$

• Parity of set of orbitals with q < p determines whether  $\hat{a}_p^{(\dagger)}$  introduces a



- BUT: we cannot represent the creation or annihilation of a particle in qubit  $q_p$ 

with  $\hat{Q}_{p}^{(\pm)}$  since  $q_{p}$  stores the parity of all orbitals with index  $q \leq p$ 

- Representing for example  $\hat{a}_{p}^{\dagger}$  in terms of  $\hat{Q}_{p}^{\pm}$  depends on qubit (p-1):
  - qubit (p-1) is in state  $|0\rangle$ : act on qubit p with  $\hat{Q}_n^+$
  - qubit (p-1) is in state  $|1\rangle$ : act on qubit p with  $\hat{Q}_n^-$



$$\hat{P}_{p}^{\pm} \equiv \hat{Q}_{p}^{\pm} \otimes |0\rangle \langle 0$$
$$= \frac{1}{2} \left( X_{p} \otimes Z_{p} \right)$$

• Operator equivalent to  $\hat{Q}^{\pm}$  in parity basis is  $\hat{P}^{\pm}$  and a **two-qubit** operator!  $0\Big|_{p-1} - \hat{Q}_p^{\mp} \otimes \Big|1\rangle \langle 1\Big|_{p-1}$  $_{-1} \mp iY_p$ 



• Operator equivalent to  $\hat{Q}^{\pm}$  in parity basis is  $\hat{P}^{\pm}$  and a **two-qubit** operator!

$$\hat{P}_{p}^{\pm} \equiv \hat{Q}_{p}^{\pm} \otimes |0\rangle \langle 0|_{p-1} - \hat{Q}_{p}^{\mp} \otimes |1\rangle \langle 1|_{p-1}$$
$$= \frac{1}{2} \left( X_{p} \otimes Z_{p-1} \mp i Y_{p} \right)$$

qubits with index greater p: apply  $\sigma_x$  to all qubits  $q_k$  with k > p

Caution: creating/annihilating a particle in p changes the parity to be stored in



• Operator equivalent to  $\hat{Q}^{\pm}$  in parity basis is  $\hat{P}^{\pm}$  and a **two-qubit** operator!

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$$= \frac{1}{2} \left( X_{p} \otimes Z_{p-1} \mp i Y_{p} \right)$$

Caution: creating/annihilating a particle in p changes the parity to be stored in qubits with index greater p: apply  $\sigma_x$  to all qubits  $q_k$  with k > p

$$\hat{a}_{p}^{\dagger} \equiv \left[\sigma_{x}^{\otimes M-p}\right] \otimes \hat{P}_{p}^{+} = \frac{1}{2} \left( \left[\sigma_{x}^{\otimes M-p}\right] \otimes X_{p} \otimes Z_{p-1} - i \left[\sigma_{x}^{\otimes M-p}\right] \otimes Y_{p} \right)$$
$$\hat{a}_{p} \equiv \left[\sigma_{x}^{\otimes M-p}\right] \otimes \hat{P}_{p}^{-} = \frac{1}{2} \left( \left[\sigma_{x}^{\otimes M-p}\right] \otimes X_{p} \otimes Z_{p-1} + i \left[\sigma_{x}^{\otimes M-p}\right] \otimes Y_{p} \right)$$



Major drawback of parity mapping: k-locality!

number of extra qubit operations which scale as  $\mathcal{O}(M)$ !

-> non-locality of "update term  $\left[\sigma_x^{\otimes M-p}\right]$ " that appears in  $\hat{a}_p^{(\dagger)}$  introduces a



Major drawback of parity mapping: k-locality!

number of extra qubit operations which scale as  $\mathcal{O}(M)$ !

• Trailing string of  $\sigma_z$  (J-W) replaced by leading string of  $\sigma_x$  ...

-> non-locality of "update term  $\left[\sigma_x^{\otimes M-p}\right]$ " that appears in  $\hat{a}_p^{(\dagger)}$  introduces a



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# Mapping the problem from fermion to qubit space

Parity mapping

fcac = np.zeros((2, 2), dtype=float) fcac[0][0] = 1.0; fcac[1][1] = 3.0

h2ac = np.zeros((2, 2, 2, 2), dtype=float) > one\_body\_ints = OneBodyElectronicIntegrals( ---

>two\_body\_ints = TwoBodyElectronicIntegrals( --fermi\_ham = qc\_ham\_qiskit.second\_q\_ops()[0] fermi\_ham.set\_truncation(0) print("Fermionic Hamiltonian",fermi\_ham) qubit\_converter = QubitConverter(mapper=ParityMapper()) number = qubit\_converter.convert(fermi\_ham) print("Parity Qubit Hamiltonian",Hq)

√ 0.4s

Fermionic Hamiltonian Fermionic Operator register length=4, number terms=5 1.0 \* ( +\_0 -\_0 )

```
+ 3.0 * ( +_1 -_1 )
+ 1.0 * (+2 -2)
+ 3.0 * ( +_3 -_3 )
+ 0j
Parity Qubit Hamiltonian 4.0 * IIII
- 0.5 * IIIZ
- 1.5 * IIZZ
 - 0.5 * IZZI
 - 1.5 * ZZII
```

```
> from qiskit_nature.converters.second_quantization import QubitConverter.
```

```
qc_ham_qiskit = ElectronicEnergy([one_body_ints, two_body_ints])
```



> from qiskit\_nature.converters.second\_quantization import QubitConverter. fcac = np.zeros((2, 2), dtype=float) fcac[0][0] = 1.0; fcac[1][1] = 3.0

h2ac = np.zeros((2, 2, 2, 2), dtype=float) > one\_body\_ints = OneBodyElectronicIntegrals( ---

>two\_body\_ints = TwoBodyElectronicIntegrals( --qc\_ham\_qiskit = ElectronicEnergy([one\_body\_ints, two\_body\_ints]) fermi\_ham = qc\_ham\_qiskit.second\_q\_ops()[0] fermi ham.set truncation(0)

# Could we avoid to to operate with $\sigma_{r}$ on all qubits q with

register length=4, number terms=5  $1.0 * (+_0 -_0)$  $+ 3.0 * (+_1 -_1)$ + 1.0 \* (+2 -2)+ 3.0 \* ( +\_3 -\_3 ) + 0j Parity Qubit Hamiltonian 4.0 \* IIII - 0.5 \* IIIZ - 1.5 \* IIZZ - 0.5 \* IZZI - 1.5 \* ZZII

q > p?



# Mapping the problem from fermion to qubit space Notes on other mappings

- to as Pauli weight
- Other more elaborate mappings exist which scale as  $O(\log_2 M)$ [Bravyi-Kitaev - > "combines ideas" of J-W and parity mappings] or even up to  $\mathcal{O}(\log_3 M)$

• Major drawback of parity and J-W mapping: k-locality which is also referred





# Mapping the problem from fermion to qubit space Bravyi-Kitaev

```
print("J-W Qubit Hamiltonian",Hq)
   qubit_converter = QubitConverter(mapper=ParityMapper())
   Hq = qubit_converter.convert(fermi_ham)
   print("Parity Qubit Hamiltonian",Hq)
   qubit_converter = QubitConverter(mapper=BravyiKitaevMapper())
   Hq = qubit_converter.convert(fermi_ham)
   print("Bravyi-Kitaev Qubit Hamiltonian",Hq)
 √ 0.7s
J-W Qubit Hamiltonian 4.0 * IIII
- 0.5 * IIIZ
- 1.5 * IIZI
- 0.5 * IZII
- 1.5 * ZIII
Parity Qubit Hamiltonian 4.0 * IIII
- 0.5 * IIIZ
- 1.5 * IIZZ
- 0.5 * IZZI
- 1.5 * ZZII
Bravyi-Kitaev Qubit Hamiltonian 4.0 * IIII
- 0.5 * IIIZ
- 1.5 * IIZZ
- 0.5 * IZII
- 1.5 * ZZZI
```

```
qubit_converter = QubitConverter(mapper=JordanWignerMapper())
Hq = qubit_converter.convert(fermi_ham)
```



# Mapping the problem from fermion to qubit space Notes on other mappings

- Major drawback of parity and J-W mapping: k-locality which is also referred to as Pauli weight
- Other more elaborate mappings exist which scale as  $O(\log_2 M)$ [Bravyi-Kitaev - > "combines ideas" of J-W and parity mappings] or even up to  $\mathcal{O}(\log_3 M)$
- Customised (tree-based) mappings tailored to QC hardware layout are also possible!





# Mapping the problem from fermion to qubit space Notes on other mappings

- Common outcome of mappings:  $\hat{H}_e$  is expressed as a sum of Pauli strings  $P_k$ 

# $_{s}a_{q} \rightarrow \hat{H}_{e} = \sum_{k} c_{k}P_{k}$



# Calculating the energy on a quantum computer

- The Hamiltonian is given as a linear combination of Pauli strings
- $\hat{H}_e = \sum c_k P_k$   $\leftarrow$  Each term is a product of local operators  $P_k = \bigotimes_{s=0}^{M-1} \sigma_{k_s,(s)}$
- We can calculate expectation values on the quantum computer

$$\left|\Psi\right\rangle = \sum_{\mathbf{f}} c_{\mathbf{f}} \left|f_{M-1} \dots f_{1} f_{0}\right\rangle \qquad \longrightarrow \qquad \left\langle\Psi\right| \hat{H}_{e} \left|\Psi\right\rangle = \sum_{k} c_{k} \left\langle\Psi\right| P_{k} \left|\Psi\right\rangle$$





# Calculating the energy on a quantum computer

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- $\hat{H}_e = \sum c_k P_k \longleftarrow$  Each term is a product of local operators  $P_k = \bigotimes_{s=0}^{M-1} \sigma_{k_s,(s)}$
- We can calculate expectation values on the quantum computer

$$|\Psi\rangle = \sum_{\mathbf{f}} c_{\mathbf{f}} |f_{M-1} \dots f_1 f_0\rangle \qquad \longrightarrow \qquad \langle \Psi | \hat{H}_e |\Psi\rangle = \sum_k c_k \langle \Psi | P_k |\Psi\rangle$$

Cannot even be written down on a classical computer



"Easy" on a quantum computer: only requires measuring Pauli strings



# Calculating the energy on a quantum computer

- The Hamiltonian is given as a linear combination of Pauli strings

$$\left|\Psi\right\rangle = \sum_{\mathbf{f}} c_{\mathbf{f}} \left| f_{M-1} \dots f_1 f_0 \right\rangle$$

Cannot even be written down on a classical computer

"Easy" on a quantum computer: only requires measuring Pauli strings



# 



# Variational optimisation with the Variational Quantum Eigensolver



# Variational optimisation with the Variational Quantum Eigensolver



... which is just one example, many algorithms exist, not to mention fault-tolerant ones...

• Prepare some quantum state using a so-called variational form (ansatz)

**Quantum Hardware** 



- Prepare some quantum state using a so-called variational form (ansatz)
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Quantum Hardware



 $|\Psi(\hat{\theta})>$ 

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• Find the ground state variationally: minimising over the parameters  $\{\vec{\theta}\}$ 



# Measuring the energy in a VQE simulation











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Every Pauli string evaluated independently through repeated measurements



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Sounds good, BUT

- How do we know the ansatz contains the ground state?
- How do we find the corresponding parameters?
- How efficient is the whole approach?
- What about the noise?



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- There are mainly two broad strategies in circuit ansatz design
  - physically motivated ansätze
  - hardware heuristic ansätze

     > parametrized circuits comprising
     single-qubit rotations and entangling
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## The Variational Quantum Eigensolver **UCCSD** ansatz

UCCSD ansatz has the form

$$\begin{split} \hat{T}_1 &= \sum_{i,a} \hat{\tau}_i^a = \sum_{i,a} \tau_i^a (a_a^{\dagger} a_i - a_i^{\dagger} a_a) \\ \left| \Psi_{\text{UCCSD}} \right\rangle &= e^{\hat{T}_1 + \hat{T}_2} \left| \Psi_{\text{ref}} \right\rangle \quad \text{with} \\ \hat{T}_2 &= \sum_{i,j,a,b} \hat{\tau}_{ij}^{ab} = \sum_{i,j,a,b} \tau_{ij}^{ab} (a_a^{\dagger} a_b^{\dagger} a_i a_j - a_j^{\dagger} a_i^{\dagger} a_b a_a) \end{split}$$

Commonly approximation: Trotterize unitary to first order lacksquare

$$\Psi_{\text{tUCCSD}} \rangle = \prod_{c \in \{ia\}} e^{\hat{\tau}_c} \prod_{d \in \{ijab\}} e^{\hat{\tau}_d} \left| \Psi_{\text{ref}} \right\rangle$$

-> Low-order trotterized form may fail to reach chemical accuracy -> Large number of exponential factors prohibits preparation on quantum processors

## The Variational Quantum Eigensolver ADAPT-VQE ansatz

- The ADAPT-VQE algorithm grows ansätze by appending one unitary at a time to a trial state
- Requires to
  - $-> \hat{\tau}^a_i$  and  $\hat{\tau}^{ab}_{ii}$  in UCCSD
  - choose a reference state calculate (HF or anything reasonable)

$$\frac{\partial E}{\partial \theta_{i}}\bigg|_{\theta_{i}=0} = \left[\frac{\partial}{\partial \theta_{i}}\left\langle \Psi^{(k)}\right| e^{-\theta_{i}P_{i}}He^{\theta_{i}P_{i}}\left|\Psi^{(k)}\right\rangle\right]\bigg|_{\theta_{i}=0} = \left\langle \Psi^{(k)}\left|[H,P_{i}]\right|\Psi^{(k)}\right\rangle$$

• The operator with the largest gradient norm is appended to the ansatz

## • a priori define an operator pool $\mathscr{P} = \{P_i\}$ , a collection of antihermitian generators

# Choosing an ansatz - ADAPT-VQE



#### Nat. Comm., 10, 3007 (2019)

# Solving the problem? The Variational Quantum Eigensolver

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# Solving the problem? The Variational Quantum Eigensolver

#### Identifying challenges towards practical quantum advantage through resource estimation: the measurement roadblock in the variational quantum eigensolver

Jérôme F. Gonthier,<sup>1</sup> Maxwell D. Radin,<sup>1</sup> Corneliu Buda,<sup>2</sup> Eric J. Doskocil,<sup>2</sup> Clena M. Abuan,<sup>3</sup> and Jhonathan Romero<sup>1</sup> <sup>1</sup>Zapata Computing, Inc., 100 Federal St., Boston, MA 02110, USA

1	Molecule	H <sub>2</sub> O	CO <sub>2</sub>	CH <sub>4</sub>	CH <sub>4</sub> O	$C_2H_6$	$C_2H_4$	$C_2H_2$	$C_2H_6O$	$C_3H_8$	$C_3H_6$	$C_3H_4$
	N <sub>el</sub>	8	16	8	14	14	12	10	20	20	18	16
	Nq	104	208	104	182	182	156	130	260	260	234	208
	$K \cdot 10^{-3}$	1.9	16	1.6	8.4	8.5	6.6	3.1	24	16	23	18
	$M \cdot 10^{-9}$	3.9	32	3.2	17	17	13	6.2	48	31	46	36
	t (days)	2.3	39	1.9	18	18	12	4.6	71	47	62	44

TABLE IV. Estimated runtimes *t* in days for a single energy evaluation using the number of measurements *M* from extrapolated

#### arXiv: 2012.04001



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# Solving the problem? The Variational Quantum Eigensolver

- How do we know the ansatz contains the ground state?
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## **Efficient measurement strategies are required! Noise mitigation is crucial!**



 $\langle E \rangle = \left\langle \Psi(\vec{\theta}) \middle| \hat{H}_e \middle| \Psi(\vec{\theta}) \right\rangle \ge E_{\text{ground}}$ 



### Molecular properties / Dynamical correlation The Variational Quantum Eigensolver Quantum Hardware

• Expectation values of one-electron  $\hat{O}_1$  and two-electron  $\hat{O}_2$  operators can be straightforward obtained from optimal VQE

$$\left\langle \hat{O}_{1} \right\rangle = \sum_{pq} o_{pq} \left\langle \hat{a}_{p}^{\dagger} \hat{a}_{q} \right\rangle$$
$$\left\langle \hat{O}_{2} \right\rangle = \sum_{pq} o_{pqrs} \left\langle \hat{a}_{p}^{\dagger} \hat{a}_{r}^{\dagger} \hat{a}_{s} \hat{a}_{q} \right\rangle$$

pqrs • Higher-order RDMs ... -> CASPT2/NEVPT2, response theory, ...

• *quantum* equation-of-motion ...



 $\langle E \rangle = \left\langle \Psi(\vec{\theta}) \middle| \hat{H}_e \middle| \Psi(\vec{\theta}) \right\rangle \ge E_{\text{ground}}$ 



# CONCLUSION

# Quantum computers without algorithms are useless machines

With proper algorithms we can make quantum computers work

We need to know how to measure

We need to mitigate errors and correct for them

